

Amendments to the Specification:

Please amend the **Detailed Description** as follows:

On page 46, please replace the paragraph found on lines 21-26 with the following paragraph:

Figure 67 illustrates an embodiment of a frequency shifter, indicated generally at 64', for use in ODN 18 of Figure 5. Frequency shifter 64' comprises a mixer 700 that is coupled to receive and shift the frequency band of RF signals in the upstream direction from diplex filter 406 for a coaxial leg 30. An output of mixer 700 is coupled through a bandpass filter 704 to combiner 408. Local oscillator 702 is coupled to provide a signal to control the operation of mixer 700.

Please replace the paragraph found beginning on page 46 line 27 to page 47 line 5 with the following paragraph:

In operation, frequency shifter 64' shifts a block of RF signals from a first frequency range to a second frequency range. For example, as mentioned above, the RF signals provided to frequency shifter 64' may comprise RF signals in the range from 5 to 40 ~~Mhz~~MHz. In one embodiment, ODN 18 comprises three frequency shifters 64'. In this embodiment, the local oscillators 702 of the three frequency shifters provide signals of 76 ~~MHZ~~MHz, 149 ~~MHZ~~MHz, and 222 ~~MHZ~~MHz, respectively. Thus, frequency shifters 64' respectively shift the upstream RF signals approximately to the 50 to 100 ~~MHZ~~MHz, 125 to 175 ~~MHZ~~MHz and 200 to 250 ~~MHZ~~MHz ranges.

On page 47, please replace the paragraph found on lines 12-16 with the following paragraph:

Figure 76 illustrates an embodiment of a frequency shifter, indicated generally at 31', for use in telephony upstream receiver 16 of Figure 83. Frequency shifter 31' returns a block of RF signals shifted by frequency shifter 64' to original frequency range of the

block. For example, frequency shifter 31' may return a block of RF signals to 5 to 40 ~~MHZ~~MHz from 50 to 100 ~~MHZ~~MHz .

Please replace the paragraph found starting on page 54 line 19 to page 55 line 9 with the following paragraph:

The present system allows a choice of different error-correction abilities for different types of data. For example, voice data is highly redundant, and needs little defense against errors. Financial transaction data, on the other hand, wants a large degree of data integrity. In addition, it may be desirable to allow a user to select -- and pay for -- whatever degree of error correction that he desires. CTSU 54, Figure 3, includes a conventional "provisioning table", which specifies a number of parameters relating to particular payload channels. Figure 55 shows a provisioning table 4411 having an added column containing indications for several different amounts of error protection. In method 4410, step 4412 reads the entry for a particular channel to be set up. In this implementation, the entry may specify message lengths of 21, 41, or 81 bits, respectively having the ability to correct 1, 2, or 4 symbols; the entry may also specify no correction, in which case message blocks do not apply. Step 4413 encodes the table entry in an IOC message and sends it to the ISU whose address appears in that row of table ~~4411~~ 4411. A general-purpose processor in CXSU 102 of the ISU stores the frame length in step 4414. As the CXSU receives data from modem 101, Figure 8, it decodes the frames of an entire message, 4415, then decodes the check symbols for the message, 4416, and signals an error, 4417, if one exists in the message. Steps 4415-4417 repeat for subsequent messages. The ISU employs the same process to send frames upstream to the head end, using the frame length setting specified in step 4414.

On page 56, please replace the paragraph found on lines 22-27 with the following paragraph:

Figure 56 shows steps ~~4120~~ 4420 for performing frame staggering. Step 4421 repeats method 4420 for all active payload channels. Step 4422 accesses the current messages for the channels in one subband. Step 4423 calculates the 1, 2, or 4 Reed-Soloman check words for the 21, 41, or 81 message data words. Step 4424 waits N frames past the start of a multiframe, whereupon step 4425 sends the message to modem 82, Figure 3 for transmission.

On page 71, please replace the paragraph found on lines 21-25 with the following paragraph:

Preferably, the MISU 66 sees approximately 3 MHz of the 6 MHz bandwidth to receive up to 130 payload channels ~~which~~ whose bandwidth also includes numerous IOC channels for communication from the HDT 12 to the MISU 66. The HISU 68 sees about 100 kHz of the 6 MHz bandwidth to receive 11 channels including at least one IOC channel for communication with the HDT 12.

Please replace the paragraph found beginning on page 72 line 19 to page 73 line 4 with the following paragraph:

Figure 58 details the operation of a typical scrambler, such as 134, Figure 21. Symbol clock 4501 clocks a seed pattern through a linear-feedback shift register 4510 having nine stages, 4510-0 through 4510-8. With XOR gate 4511 positioned as shown, the generator polynomial is binary "100 010 000". The seed initially loaded into register 4510 at input 4502 is "111 001 110". Two identical translation tables 4520 and 4521 receive two-bit inputs from register 4510 at every symbol time. The high- and low-order bits of table 4520 proceed from the outputs of stages 4510-~~8~~7 and 4510-~~7~~6, respectively. High-order bit 4523 of table 4521 also receives output 4510-~~7~~6, but as its high-order bit; stage output 4510-~~6~~5 provides its low-order bit. Logic gate[[s]] 4530 perform an XOR between the five-bit output of table 4520 and the upper five bits of a 10-bit DS0 word, while gate[[s]] 4531 does the same for the

five-bit output of table 4521 and the lower five bits of the same DS0 word. Outputs 4505 and 4506 carry the two 5-bit scrambled symbols for the DS0 word. Each descrambler such as 176, Figure 22 or 23, is identical to its corresponding scrambler. It recovers the original bit pattern of each symbol by decoding it with the same polynomial and seed.

On page 74, please replace the paragraph found on lines 10-24 with the following paragraph:

Within the framework of QAM32 modulation, Figure 17 shows a constellation which has improved characteristics. Here, the in-phase and quadrature values are shown encoded by three bits each instead of the four shown in Figure 14+3; ~~their~~ analog values, however, ~~they~~ remain in the ranges -5 to +5. The constellation of Figure 17 approaches as closely as possible to an analogy to a Gray code scheme, in which a transition from one row to the next and from one column to the next result in only a single bit change in the 5-bit symbol code. (The exceptions are four transitions from the first column to the second, and from the fifth to the sixth, which have two transitions each. The corner cells have zero transitions between these columns, which do not detract from the advantages of the scheme.) If a symbol is received incorrectly after transmission, the most likely error is a slight change in either amplitude or phase. If the bit strings represented by the symbols have as few bit transitions as possible for single-value phase and amplitude changes, then a reception error will create fewer bit errors on the final digital output. That is, small (symbol) errors in produce small (bit) errors out.

On page 75, please replace the paragraph found on lines 13-25 with the following paragraph:

Each symbol that gets represented by the I/Q values is mapped into a fast Fourier transform (FFT) bin of symbol buffer 138 in Figure 21. For example, for a DS0+, running at 8 kHz frame rate, five bits are mapped into one FFT bin and five bits into another bin. Each bin or memory location of the symbol buffer 138 represents the payload data and control data in the frequency domain as I/Q values. One set of FFT bins gets mapped into the time domain through the inverse FFT 140, as is known to one skilled in the art. The inverse FFT 140 maps the complex I/Q values into time domain samples corresponding to the number of points in the FFT. Both the payload data and IOC data are mapped into the buffer 138 and transformed into time domain samples by the inverse FFT 140. The number of points in the inverse FFT 140 may vary, but in the preferred embodiment the number of points is 256. The output of the inverse FFT 140, for a 256 point FFT, is 256 time domain samples of the waveform.

On page 76, please replace the paragraph found on lines 9-25 with the following paragraph:

In a presently preferred embodiment, the 256 samples clocked into inverse FFT 140 represent an extra 45° ($\pi/4$ radians) above a complete cycle. Another way to think of this is that the symbols are clocked into the FFT at an effective 9 kHz rate, and clocked out at the nominal 8 kHz symbol rate. Figure 52 shows an unmodulated sine wave, *i.e.*, one having $I=0$, $Q=0$ in the units used herein. The upper portion shows one cycle, 0-360°, at the nominal 8 kHz frame rate. The lower portion shows the same wave at a 9 kHz rate, so that the amount of time previously occupied by 360° now takes up 405° of phase -- from -22.5° to +382.5°. Obviously, there are phase discontinuities between successive cycles of the wave. Figure 53 shows a typical QAM 32 wave modulated at a different amplitude and a slightly different phase from those of Figure 52. These might correspond to, say, $I=-1$, $Q=+1$ in the scheme used herein. The small portions at the ends of this wave represent unmodulated cycles, as in Figure 52. The phase of this wave is advanced from the corresponding wave of the lower portion of Figure 52; it does not

cross the zero axis at 0° and 180° of its proper cycle. It does, however, include the extra 22.5° of excess phase at each end, for 45° extra over an 8 kHz cycle. Again, [[a]] phase discontinuities exist at the ends of the total 405° phase degrees of this wave.

On page 79, please replace the paragraph found on lines 15-24 with the following paragraph:

RF master clock ~~42514261~~ proceeds to RF synthesizer 143 in HDT transmitting modem 82, as shown in Figure 21, where it directly controls the frequency of the tunable 500-850MHz RF carrier for the entire band carrying all of the channels shown in Figures 13 and 16. Symbol clock 4261 proceeds to the frame-clock inputs in Figure 21, where it controls the symbol timing, and, because it also controls the FFT speed, the frequencies of the channels in the entire band. Clock lock 4200 thus provides a solid link which inherently preserves the orthogonality of the band signals in a multicarrier system, by deriving the RF carrier clock and the symbol or frame clock from the same source. At the same time, it provides a small amount of gradual variation for satisfying the demands of the analog RF components.

Please replace the paragraph found beginning on page 79 at line 25 to page 8 line 5 with the following paragraph:

The overall purpose of locking the two clocks together at the HDT is to lock the carrier clocks and the symbol (frame) clocks throughout the system; and the purpose of this in turn is to preserve the orthogonality of the signals in a multicarrier system which is capable of bidirectional operation: that is, as a multipoint-to-point-configuration as well as in the usual point-to-multipoint "broadcast" direction. Clock generator 166, Figures 22 and 23, of timing generator 107, Figure 68 locks to the frequencies of the incoming signals to provide the clocks used in the remote ISU modules. Therefore, the

carrier and frame clocks in each upstream transmitter portion, Figure 24, of remote modem 108, Figure 8, are also locked to each other, by virtue of being locked to the incoming signal from the HDT.

On page 106, please replace the paragraph found on lines 5-24 with the following paragraph:

Simultaneously with symbol alignment, the CXMC 80 transmits a message to the MCC modem 82 to perform path delay adjustment. The CXMC 80 sends a message on a downstream IOC channel via the MCC modem 82 instructing the CXSU controller 102 to enable the ISU modem 101 to transmit $[[\star]]$ another signal on a synchronization channel which indicates the multiframe (2 kHz) alignment of the ISU 100. The MCC modem 82 detects this multiframe alignment pattern and performs a hardware correlation on the pattern. From this correlation, the modem 82 calculates the additional symbol periods required to meet the round trip path delay of the communication system. The MCC modem 82 then returns a message to the CXMC 80 indicating the additional amount of delay which must be added to meet the overall path delay requirements and the CXMC then transmits a message on a downstream IOC channel via the MCC modem 82 instructing the CXSU controller 102 to relay a message to the ISU modem 101 containing the path delay adjustment value. Numerous iterations may be necessary to reach path delay equilibrium and if it is not reached within a predetermined number of iterations, then the ISU may again be reset. Such adjustment is made in the ISU transmitter as can be seen in the display delay buffer "n" samples 192 of the upstream transmitter architectures of Figures 24 and 25. Path delay and symbol alignment may be performed at the same time, separately or together using the same or different signals sent on the synchronization channel.

On page 107, please replace the paragraph found on lines 1-16 with the following paragraph:

After an ISU 100 is initialized and activated for the system, follow-up synchronization or tracking may be performed periodically to keep the ISUs calibrated within the required tolerance of the OFDM transport requirements. The follow-up process is implemented to account for drift of component values with temperature. If an ISU 100 is inactive for extreme periods of time, the ISU can be tuned to the synchronization channels and requested to update upstream synchronization parameters in accordance with the upstream synchronization process described above. Alternatively, if an ISU has been used recently, the follow-up synchronization or tracking can proceed over an IOC channel. Under this scenario, as generally shown in Figure 28, the ISU 100 is requested to provide a signal over an IOC channel by the HDT 12, 2800. The HDT 12 then acquires and verifies that the signal is within the tolerance required for a channel within the OFDM waveform 2811. If not, then the ISU is requested to adjust such errored parameters 2813. In addition, during long periods of use, ISUs can also be requested by the HDT 12 to send a signal on an IOC channel or a synchronization channel for the purpose of updating the upstream synchronization parameters.

On page 118, please replace the paragraph found on lines 3-14 with the following paragraph:

The steps in the left column of Figure 50 are performed within each ISU; the HDT performs the steps on the right. Step 4110 selects a number of payload channels for monitoring from table 4111. The channels must include one channel from each separate ISU, but need not include more than one. An MISU thus needs time in only one of its 130 payload channels, a very low overhead. A 10-channel HISU subband, however, may require time in more than one of its 10 payload channels, because multiple HISUs can share the same subband. Of course, a powered-down ISU, or one having no upstream payload channels provisioned to it, need not participate in blocks ~~4740~~3940 of Figure 47, because it does not transmit upstream at all. (It is alternatively possible to employ IOC channels instead of payload channels for this purpose. Although requiring less overhead, such use is generally much more complex to implement.)

On page 118, please replace the paragraph found on lines 15-25 with the following paragraph:

After ranging procedure 3900 has acquired a correct initial power level, step 4120 performs a scan every 30msec. for all the selected payload channels, as indicated by arrow 4121. Each ISU responds at 4130 by sending a message 4131 on its selected upstream channel when the scan reaches that channel. In step 4140, the HDT measures the received power level from each ISU separately. If the signal amplitude is within a certain range of its previous value, then steps 4150 compensate for the variation at the HDT. Step 4151 smooths the errors over several scans, to prevent sudden jumps from a single glitch. Step 4152 then adjusts the coefficients of equalizer 214 in the upstream receiving portion, Figure 26, of modem 82, Figure 3, to compensate for the variation. This sequence compensates for small, slow variations at a high resolution; the equalizer steps are small and very linear.

On page 119, please replace the paragraph found on lines 23-28 with the following paragraph:

Thus, power-leveling blocks ~~4740~~3950 take advantage of the characteristics of the system to adjust both ends in a way which achieves both high resolution and wide dynamic range. The digital control available in the head-end equalizer provides precision and linearity in tracking slow changes, and the analog control at the ISU provides a wide range, and still allows the head end to track out inaccuracies caused by its coarse and nonlinear nature.

Please replace the paragraph found beginning on page 124 line 29 to page 125 line 8 with the following paragraph:

Referring to Figures 41, 42, and 43, the basic short integration operation is described. When a parity error 5000 of a channel is detected by the CXMU 56, a parity interrupt is disabled by setting the interrupt priority level above that of the parity

interrupt 5001 (Figure 41). If a modem alarm is received which indicates a received signal failure, parity errors will be ignored until the failure condition ends 5002. Thus, some failure conditions will supersede parity error monitoring. Such alarm conditions may include loss of signal, modem failure, and loss of synchronization. If a modem alarm is not active 5004, a parity count table is updated 5006 and an error timer event as shown in Figure 42 is enabled.

On page 125, please replace the paragraph found on lines 9-18 with the following paragraph:

When the error timer event is enabled 5100, the channel monitor 296 enters a mode wherein parity error registers of the CXMU 56 are read every 10 milliseconds and error counts are summarized after a one second monitoring period 5105. Generally, the error counts 1240 are used to update the channel quality database 5334 and determine which (if any) channels require re-allocation. The channel quality table 300 of the database contains an ongoing record of each channel. The table organizes the history of the channels in categories such as: current ISU assigned to the channel, start of monitoring, end of monitoring, total error, errors in last day, in last week and in last 30 days, number of seconds since last error, severe errors in last day, in last week and in last 30 days, and current service type, such as ISDN, assigned to the channel.

Please replace the paragraph found beginning on page 125 line 19 to page 126 line 3 with the following paragraph:

As indicated in Figure 41, after the parity interrupt is disabled and no active alarm exists, the parity counts are updated 5006 and the timer event is enabled 5008. The timer event (Figure 42), as indicated above, includes a one second loop where the errors are monitored. As shown in Figure 42, if the one second loop has not elapsed 5110, the error counts are continued to be updated 5104. When the second has elapsed 5106, the errors are summarized 5120. If the summarized errors over the one second period exceed an allowed amount indicating that an allocated channel is corrupted or bad 5121, as

described below, channel allocator 304 is notified 5123 and ISU transmission is reallocated to a different channel. As shown in Figure 43, when the reallocation has been completed 5200, the interrupt priority is lowered below parity 5210 so that channel monitoring continues and the channel quality database is updated 5215 concerning the actions taken. The reallocation task may be accomplished as a separate task from the error timer task or performed in conjunction with that task. For example, the reallocator 304 may be part of channel monitor 296.

On page 126, please replace the paragraph found on lines 4-8 with the following paragraph:

As shown in Figure 44 in an alternate embodiment of the error timer task 5110-2 of Figure 42, channels can be determined to be bad 5304 before the one second has elapsed. This allows the channels that are determined to be corrupted during the initial portion of a one second interval to be quickly identified and reallocated 5308 without waiting for the entire one second to elapse.

On page 126, please replace the paragraph found on lines 9-18 with the following paragraph:

Instead of reallocation, the power level for transmission by the ISU may be increased to overcome the ingress on the channel as indicated in Figure 124. However, if the power level on one channel is increased, the power level of at least one other channel must be decreased as the overall power level must be kept substantially constant. If all channels are determined bad 5306, the fault isolator 302 is notified 5320 indicating the probability that a critical failure is present, such as a fiber break. If the summarized errors over the one second period do not exceed an allowed amount indicating that the allocated channel is not corrupted, the interrupt priority is lowered below parity 5210 and the error timer event is disabled 5332. Such event is then ended 5350 and the channels once again are monitored for parity errors per Figure 41.

Please replace the paragraph found beginning on page 126 line 19 to page 127 line 3 with the following paragraph:

Two issues presented by periodic parity monitoring as described above must be addressed in order to estimate the bit error rate corresponding to the observed count of parity errors 1240 in a monitoring period of one second to determine if a channel is corrupted. The first is the nature of parity itself. Accepted practice for data formats using block error detection assumes that ~~an errored a block error~~ represents one bit of error, even though the error actually represents a large number of data bits. Due to the nature of the data transport system, errors injected into modulated data are expected to randomize the data. This means that the average ~~errored block error~~ frame will consist of four (4) ~~errored block error~~ data bits (excluding the ninth bit). Since parity detects only odd bit errors, half of all ~~errored block error~~ frames are not detected by parity. Therefore, each parity (frame) error induced by transport interference represents an average of 8 (data) bits of error. Second, each monitoring parity error represents 80 frames of data (10 ms/125 μ s). Since the parity error is latched, all errors will be detected, but multiple errors will be detected as one error.

On page 129, please replace the paragraph found on lines 8-16 with the following paragraph:

The following is a description of the long integration operation performed by the background monitor routine (Figure 45) of the channel monitor 296. The background monitor routine is used to ensure quality integrity for channels requiring greater quality than the short integration 10^{-3} bit error rate. As the flow diagram shows in Figure 45, the background monitor routine operates over a specified time for each service type, updates the channel quality database 6006 table 300, clears the background count 6008, determines if the integrated errors exceed the allowable limits determined for each service type 6010, and notifies the channel allocator 304 of bad channels as needed 6012.

On page 129, please replace the paragraph found on lines 17-26 with the following paragraph:

In operation, on one second intervals, the background monitor updates the channel quality database 6006 table. Updating the channel quality data table has two purposes. The first purpose is to adjust the bit error rate and number of errored seconds data of error free channels to reflect their increasing quality. The second purpose is to integrate intermittent errors on monitored channels which are experiencing error levels too low to result in short integration time reallocation (less than 4 parity errors/second). Channels in this category have their BER and numbers of errored seconds data adjusted, and based on the data, may be re-allocated. This is known as long integration time re-allocation, and the default criteria for long integration time re-allocation for each service type are shown as follows:

Please replace the paragraph found beginning on Page 130 line 20 to page 131 line 6 with the following paragraph:

Unallocated or unused channels, but initialized and activated, whether used for reallocation for non-concentration services such as TR-8 or used for allocation or reallocation for concentration services such as TR-303, must also be monitored to insure that they are not bad, thereby reducing the chance that a bad channel will be allocated or reallocated to an ISU 100. To monitor unallocated channels, channel monitor 296 uses a backup manager routine (Figure 46) to set up unallocated channels in a loop in order to accumulate error data used to make allocation or re-allocation decisions. When an unallocated channel experiences errors 6110, it will not be allocated to an ISU 100 for one hour 6118. After the channel has remained idle (unallocated) for one hour, the channel monitor places the channel in a loop back mode 6120 to see if the channel has improved. In loop back mode, the CXMU 56 commands an initialized and activated ISU 100 to transmit a message on the channel long enough to perform short or long integration on the parity errors as appropriate. In the loop back mode, it can be determined whether the previously

corrupted channel has improved over time and the channel quality database is updated accordingly. When not in the loop back mode, such channels can be powered down.

Please replace the paragraph found beginning on page 161 line 28 to 162 line 10 with the following paragraph:

The Sigma Delta decimator system (SDDS) 2850 is a N-bit A-to-D converter which generates a one-bit serial data stream having resolution and accuracy of N bits (in one embodiment, 15-bit resolution is obtained; in another embodiment, the A/D has a 10-bit resolution with 9-bit linearity). SDDS 2850 is running on the clock generator 2749 divided to 73.728 MHz which oversamples the SDDS input signal in order that it only passes data at the 18.432 MHz \pm 100 kHz, approximately. The following circuits 2841 - 2847 then take that 200 kHz of frequency that Sigma-Delta converter 2840 passes and shifts it down a base band, ~~so basically it goes~~ providing a range from 0 Hz to -200 kHz. This a SDDS 2850 turns this relatively slow signal ~~and then SDDS 2850 turns that~~ into 10-bit parallel words. The Sigma-Delta converter 2840 ~~is outputting~~ outputs a 1-bit serial ~~1-bit~~ stream, which is ANDed with two 18.432 MHz square waves to produce serial digital I and Q that are two 18.432 MHz gated square waves.

On page 166, please replace the paragraph found on lines 10-16 with the following paragraph:

The system 500 will utilize the telephony error correction mechanism described above with respect to system 10 of Figure 1. Under the telephony error correction scheme, forward error correction codes are generated for upstream traffic but not for downstream traffic. Forward error correction on upstream is generated at the ISU 100 (HISU 68 or MISU 66) and consumes the 10th bit of each DS0, thereby protecting each DS0 separately. The error correction can be disabled, but this is not recommended for data transport.

On page 166, please replace the paragraph found on lines 17-22 with the following paragraph:

The error detection/correction processing occurs on the CXMU 56 of Figure 3 and corrected data is delivered ~~corrected~~ to LANUs 580, DS1Us 48, in MARIO streams. Therefore, the system 500 data architecture does not explicitly have to deal with error correction. The CRC of the HDLC frames provide for a level of error detection above the error detection/correction of the CXMU 56. Errors detected in the LANU 580 will be reported through the SNMP agent of the LANU 580.

On page 173, please replace the paragraph found on lines 7-19 with the following paragraph:

For upstream traffic, the HDLC framed data available on the MARIO interface passes through the rate adaptation and TSA block ~~488~~588. In this block, the 2.56 Mbps MARIO interface is rate adapted down to 2.048 Mbps. As part of rate adaption, the ninth bit of each DS0+ of the MARIO stream is stripped and sent to the NBS logic. The ninth bit carried an order number that is used to time order the DS0s in multi-channel calls. Once the order numbers are established, the processor 581 configures the TSA to re-order the multi-channel calls and target the super-channel to a QMC 586. For super-channels composed of 32 or less DS0s, the call is placed in a single 2.048 Mbps data stream, along with other calls and sent to a QMC 586. For 128 DS0 calls, the DS0s are placed in a single 8.192 Mbps stream that is target to a QUICC32 HDLC controller 586 configured for 8.192 Mbps HDLC. Whether targeted to a 2.048 Mbps QMC or 8.192 Mbps HDLC, the frames are accumulated and transmitted on the local Ethernet LAN.

Please replace the paragraph found beginning on page 173 line 21 to page 174 line 5 with the following paragraph:

For downstream traffic, data on the LAN is filtered according to the destination MAC address. If the MAC address is in the Content Addressable Memory (CAM) 589, the LANU 580 will accept the Ethernet frame. Once the frame is accepted, the LANU

~~588~~580 accesses a routing table in memory 582 to select the appropriate MARIO slot for transport. The frame is then scheduled for transmission and the HDLC controller 586 takes over. In the downstream direction for 32 DS0 or less calls, the HDLC controller 586 is responsible for creating the MARIO stream and encoding the data into HDLC format. For 128 DS0 calls, the 8.192 Mbps HDLC data stream is split among the four MARIO interfaces (A-D). The time ordering and aggregation in the downstream direction is controlled by the TSA. After the data passes through the TSA, the ninth bit signaling information is added to indicate that data transmission is enabled. At the same time that the ninth bit signal is added, the data stream is rate adapted up to the 2.56 Mbps rate of the MARIOs.

On page 176, please replace the paragraph found on lines 18-27 with the following paragraph:

The basic design of the DMCU 560 is very similar to the design of DMSM 550. The most notable difference is the interface to the RF portion is two 8.192 Mhz serial channels. This allows the MISU interface to support a symmetrical 8.192 Mbps Ethernet connection. Because of the higher throughput the MISU is based on the MC68360 614 that can support both the 8.192 Mbps HDLC connection as well as the 10 Mbit Ethernet. In front of the 10 Mbit Ethernet interface 612 is a router 616 that allows four users access to the 8.192 Mbit data connection. The router 616 design ~~insures~~ensures security for all connected users. The design contains 2 Mbytes of DRAM 618 and 256bytes of FLASH 619. Like the HISU design, the FLASH can be remotely updated with TFTP.

Please replace the paragraph found beginning on page 176 line 28 to page 177 line 12 with the following paragraph:

The DMCU 610 has an equivalent interface function that moves data from the HDLC controller 611 to the RF modem and works in a very similar way to the DMSM 600 with the exception that the MISU modem interface is formatted as two SLUIGI streams that are clocked at 8.192 Mbps. Between the two SLUIGI streams, 128 DS0s

can be carried between the HDLC controller 611 and the modem. The MISU interface logic 613 is responsible for buffering and then sending data over the dual S-LUIGI interface to and from the RF modem. In the upstream direction the MISU interface logic 613 generates order numbers for each of the 128 DS0s over the ninth bit of the DS0+. ~~The +~~ (the order numbers generated on the MISU work in the same way as they do on the HISU). In the downstream direction, the MISU interface logic 613 is responsible for moving data from the S-LUIGI streams to the HDLC controller 611. The MISU interface logic also monitors the ninth bit data stream from the first DS0 to detect the "Data Dial Tone" that enables data transmission. Ethernet controller 617 moves the data to the router 616.

On page 180, please replace the paragraph found on lines 7-12 with the following paragraph:

Also, using the ninth bit signaling, the mLANU can "steal" bandwidth from other (e.g. high capacity) users. This is done by removing the Data Dial Tone from a subscriber using the ninth bit signaling. This quiesces the user's line, allowing, the number of channels for that user can be reassigned to increase the bandwidth allocated to the user. This technique is can also be used to decrease the number of channels assigned to a user.

On page 180, please replace the paragraph found on lines 13-20 with the following paragraph:

In order to establish the transport, the CXMU 56 trains the modems (as described above with respect to system 10) and associates the available tones with DS0s. Once the training is complete, the CXMU 56 sends a "Pass" message to the CTSU 54, which in turn informs the mLANU over the TMC that the call is complete with the "Call Complete" message. In response, the mLANU ~~580~~ configures the HDLC controllers 586 and the TSA 588 on the mLANU or another LANU through communications over the backplane LAN. At this point the pipe is established but data is not yet enabled.

On page 181, please replace the paragraph found on lines 2-16 with the following paragraph:

A session is terminated at the customer end when no data is available for transmission by generating an "On Hook" message. The call processing sequence for an "On Hook" message is shown in Figure 110. When the CDM terminates the connection, an "On Hook" message is sent over the IOC to the CXMU 56. The CXMU 56 ~~is in~~ response sends an "On Hook" message, identifying the CRV, to the CTSU 54. The CTSU 54 then sends a "Tear Down" message to the mLANU 584 over the TMC. At the mLANU 580, the connection is deleted from the connection database and then released. If the connection is not on the mLANU, the mLANU will send a "Release Channel" message to the target LANU 580 and also will send a "Release Cross Connect" message to the CTSU 54. The CTSU 54 will release the cross connects used for the connection and then send a "Release Bandwidth" message to the CXMU 56. At the CXMU 56 the mapping between tones and DS0s is lost and the connection is lost. When the connection is lost, the CDM will lose the "Data Dial Tone" in the ninth bit signaling of the first DS0 of the call.

On page 182, please replace the heading and paragraph found on lines 1-10 with the following heading and paragraph:

LANU ~~580~~ Software

The LANU software 620 is responsible for the all major function of the data concentration of the head-end equipment. A simplified schematic diagram of the LANU software is shown in Figure 111. The ~~LANU software 620 of the LANU~~ consists of three major components: bridging 621, HDLC LAN manager 622, and data IDT 623. All three tasks will operate as applications on top of the embedded controller operating system "pSOS" kernel in the processor 581. The pSOS kernel will provide the base for the multi-tasking operation of the software 620.

On page 182, please replace the paragraph found on lines 11-21 with the following paragraph:

The most important task to the actual transport of data will be the bridging task. The bridging task has several functions. First, the task will be responsible for providing the virtual switch between the MARIO and Ethernet interfaces. The task will be implemented as ~~in an~~ "interrupt on receive" task that will execute at interrupt level. At either interface, an interrupt is issued when an entire frame has been received and stored in buffer memory (Figure 103, 582). During the interrupt service routine, the packet will be handed off by modifying the associated buffer descriptor after looking up the routing in the bridging table (stored in memory 582). For upstream traffic (HDLC to Ethernet), the source of the first packet will be read to discover its MAC address. This address will be added to the bridging table and written to the Ethernet CAM 589 for filtering.

On page 194, please replace the paragraph found on lines 9-25 with the following paragraph:

In each case in which an ISU is assigned to a subband, channel manager 900 uses various criteria to select the subband for an ISU. Figure 62 is a flow chart that illustrates one embodiment of a method for assigning an ISU to a subband. According to this method, channel manager 900 first selects a subband 6202. Channel manager 900 then determines whether addition of the ISU to the subband would provide an acceptable load on the IOC channel 6204. For example, channel manager considers the number of ISUs assigned to a subband. Further, channel manager considers the type of ISU and the likely load that the ISU will place on the IOC channel. By considering these factors, channel manager 900 can selectively distribute the load on the IOC channels so as to facilitate timely communication of control data to and from the ISU. This also allows channel manager 900 to evenly distribute the ISUs over the available subbands such that a like number of ISUs occupy each subband. Channel manager 900 also weighs the number of available channels 6206 within the subband and their transmission quality 6208 as

recorded in the tables of channel manager 900. Channels with longer low-error rate histories will be used first. Channels previously marked bad and reallocated for monitoring will be used last. Based on these criteria, channel manager selects a subband for each ISU 6210.

Please replace the paragraph found beginning on page 195 line 22 to page 196 line 17 with the following paragraph:

Channel manager 900 receives a request for allocation of a payload channel from either the SCNU 58 or CTSU 54. At block 912, channel manager 900 decides whether sufficient payload channels are available in the current subband to fulfill the request. If sufficient channels are available, the method proceeds to block 914 and determines whether one of the available channels is the IDL channel. If the IDL channel is not one of the available channels, channel manager 900 allocates a channel for each channel requested by CTSU 54 or SCNU 58 at blocks 916 and 918. Channel manager 900 selects the channels based on several criteria that increase the likelihood of achieving a connection with acceptable quality levels. For example, channel manager 900 can use the method shown in Figure 61. According to this method, channel manager 900 begins the selection process 6100 by identifying available payload channels 6102 that are located toward the center of the 6 MHz transmission channel. Typically, channels that are nearer to the edge of the 6 MHz channel exhibit higher bit error rates than the channels that are closer to the center. Further, channel manager 900 can also consider limitations of the ISU and the requested service in selecting a payload channel. For example, the ISU may be preset for use only with odd or even payload channels. This information may be included in a ROM on the ISU and provided to the channel manager when channel allocation is requested or during acquisition. Further, channel manager 900 uses data on the quality of transmissions over the identified channels stored in tables in channel manager 900 to determine which available payload channels have an acceptable error history 6104, e.g., bit error rate. Other appropriate criteria can be used in channel selection that also tend to increase the chances of producing a connection with

acceptable quality 6104. Based on these criteria, channel manager selects a payload channel to allocate to the ISU 6106.

On page 199, please replace the paragraph found on lines 6-22 with the following paragraph:

Figure 68 is a flow chart that illustrates a method for monitoring payload channels by channel manager 900. Channel manager 900 reads parity error registers 901 of the CXMU 56 are read every 10 milliseconds. Generally, the error counts are used to update the channel quality database and determine which (if any) channels require reallocation. The database of channel manager 900 contains an ongoing record of each channel 903. An accumulator sums the errors 905 with previously recorded errors to update the database. The database organizes the history of the channels in categories such as: current ISU assigned to the channel, start of monitoring, end of monitoring, total error, errors in last day, in last week, number of seconds since last error, severe errors in last day, in last week, and current service type, such as ISDN, assigned to the channel. When the channel is a regular (non-loop back) payload channel 907, channel manager 900 determines whether the performance statistics in the database are within service specific threshold 909. When the statistics unacceptably exceed the threshold 910, channel manager 900 reallocates the channel 911 using a "make before break" procedure to reduce the disruption from reallocating the channel. Thus, channel manager 900 allocates the new payload channel for the connection before deallocating the current payload channel.

Please replace the paragraph found beginning on page 203 line 22 to page 204 line 4 with the following paragraph:

Figure 69 is a flow chart that illustrates an embodiment of a method for allocating a payload channel to the ISU data link. At block 330a, channel manager 900 receives a request for an IDL channel. At block 332a, channel manager 900 determines whether a payload channel is available. If a payload channel is available, channel manager 900

allocates the payload channel to the IDL channel 335 and the data is transmitted to the identified ISU. If, however, a channel is not available, channel manager determines whether one of the allocated channels is idle by checking the hook state of a line of a channel unit 342. If the line is on hook 339, then channel manager 900 reallocates the channel to the IDL channel 343 until the IDL transmission is complete. If however, channel manager receives a request for a communication link to the line of the channel unit, channel manager interrupts the IDL channel and reallocates the payload channel to the channel unit.

On page 204, please replace the paragraph found on lines 18-23 with the following paragraph:

When channel manager 900 determines that an ISU can be powered down, the ISU's transmitter is disabled. Head end 32 detects the loss of power to the ISU and sends an idle pattern upstream to the switch. An IOC control message to ~~ofor~~ from the IOC will power-up the ISU. The IOC traffic to or from the ISU indicates to the background task in charge of powering down ISUs to check the ISU against the criteria for powering down again.

On page 206, please replace the paragraph found on lines 2-15 with the following paragraph:

The present invention shall now be described in further detail with reference to Figures 116-123. The first part of the description shall primarily deal with downstream transmission and the second part of the description shall primarily be with regard to upstream transmission. The video and telephony distribution network 1006 in accordance with the present invention, includes head end 1010 which receives video and telephony information from video and telephony service providers via trunk line 1008. Head end 1010 includes a plurality of host distribution terminals 1012 and a video host distribution terminal 1014. The HDT 1012 includes ~~[[a]]~~ transmitters and receivers for communicating telephony information, such as T1, ISDN, or other data services

information, to and from telephony service providers via trunk line 1008 and the VHDT 1014 includes [[a]] transmitters and receivers for communicating video information, such as cable TV video information and interactive data of subscribers to and from video service providers via trunk line 1008.

On page 213, please replace the paragraph found on lines 9-24 with the following paragraph:

In another embodiment of providing downstream electrical video and telephony signals from the coaxial taps 1034 to remote units 1042, as shown in Figure 116, a separate coaxial line ~~from~~from coaxial tap 1034 is utilized to provide transmission of the signals therefrom to set top box 1078, and thus for providing the downstream video signals to video equipment unit 1072. In such a case, a second coaxial line from coaxial tap 1034 would be utilized to provide the downstream telephony signals to a multiple integrated service unit (MISU) 1044 which would be much like the home integrated service unit 1070 as described with regard to Figure 119 except lacking an ingress filter 1098 and tap 1097. Unlike home integrated service unit 1070, the MISU 1044 would be utilized to service several remote units 1042 with telephony services via various service interfaces. Whether the video and telephony signals are provided to the curb with use of the MISU 1044 or whether the video and telephony signals are provided directly to a home integrated service unit is strictly one of application and either can be utilized with regard to the same or different coaxial taps 1034 and within the same or different coaxial distribution systems 1007.

On page 214, please replace the paragraph found on lines 5-26 with the following paragraph:

The above description primarily involves the downstream transmission of video and telephony information from head end 1010 to remote units 1042. The upstream transmission of interactive data from set top boxes 1078 and other data, for example telephony from telephones 1076, shall now be described with reference to Figures 116-

123. The description shall be limited to transmission from remote units via home integrated service units as transmission from an MISU is substantially similar and easily ascertainable from such description. Home integrated service unit 1074 provides set top box information from set top box 1078 and telephony information from the service interfaces 1112, including information from telephone 1076, to the optical distribution ~~modem~~ node 1026 connected thereto by the same coaxial path as for the downstream communication. The set top box signals are transmitted by a separate RF modem of the video service provider at a relatively low frequency in the bandwidth of about 5 to 40 MHz which is unused by telephony and video services. The telephony signals are also transmitted upstream in the 5-40 MHz bandwidth, usually from 10 MHz to 30 MHz. This 5-40 MHz bandwidth is reused in each coaxial path 1029 from each remote unit 1042 to the respectively connected optical distribution node 1026. As such, upstream electrical telephony data signals from the remote units are transmitted at the same reused frequency bandwidth of 5-40 MHz on each coaxial line 1029 for input to the optical distribution node 1026. Therefore, as shown in Figure 118, four upstream electrical telephony signals, each in the 5-40 MHz bandwidth, are input to optical distribution node 1026, via the respectively connected coaxial cables 1029.